

This question paper consists of 8 printed pages, each of which is identified by the Code Number COMP5450M.

A non-programmable calculator may be used.  
Answer All Questions.  
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School of Computing

**January 2019**

**COMP5450M**

KNOWLEDGE REPRESENTATION AND REASONING (MSc)

Time allowed: 2 hours

**PLEASE DO NOT REMOVE THIS PAPER FROM THE EXAM ROOM**

**Answer ALL THREE questions**

The marks available for each part of each question are clearly indicated.

## Question 1

(a) Translate the following sentence into *Propositional Logic*:

- I go shopping on Mondays and Tuesdays.

[2 marks]

**Answer:**  $(Monday \vee Tuesday) \rightarrow Shopping.$

(b) Translate the following sentences into *First-Order Predicate Logic* (using equality where necessary):

- (i) Some yellow frogs are poisonous.

[2 marks]

**Answer:**  $\exists x[Frog(x) \wedge Yellow(x) \wedge Poisonous(x)]$

- (ii) All Helen's rabbits are white or grey.

[2 marks]

**Answer:**  $\forall x[(Rabbit(x) \wedge Owns(helen, x)) \rightarrow (White(x) \vee Grey(x))]$

- (iii) No dog ate more than one biscuit.

[2 marks]

**Answer:**  $\neg \exists x \exists y \exists z [Dog(x) \wedge Biscuit(y) \wedge Biscuit(z) \wedge \neg(y = z) \wedge Ate(x, y) \wedge Ate(x, z)]$

- (iv) Edward hates everyone except himself.

[2 marks]

**Answer:**  $\forall x[Hates(ed, x) \leftrightarrow \neg(x = ed)]$

(c)  $\mathcal{M} = \langle \mathcal{D}, \delta \rangle$  is a model for a first-order language with two unary predicates  $P$  and  $Q$  and a binary relation predicate  $R$ . The domain of  $\mathcal{M}$  is the set  $\{a, b, c, d, e, f\}$ , and the denotation of the predicates is:

- $\delta(P) = \{a, b, c\}$
- $\delta(Q) = \{d, e, f\}$
- $\delta(R) = \{\langle a, f \rangle, \langle b, e \rangle, \langle c, d \rangle, \langle f, f \rangle\}$

Which of the following formulae are satisfied by this model?

[4 marks]

F1.  $\forall x[P(x) \vee Q(x)]$

F2.  $\exists w[P(w) \wedge Q(w)]$

F3.  $\forall x[P(x) \rightarrow \exists y[R(x, y) \wedge Q(y)]]$

F4.  $\neg \exists x \exists y[Q(x) \wedge Q(y) \wedge R(x, y)]$

**Answer:** F1 yes, F2 no, F3 yes, F4 no. 1 mark each.

(d) Use the *Sequent Calculus* to show that the following sequent is valid: **[6 marks]**

$$\forall x[P(x)], \forall x[P(x) \rightarrow Q(x)] \vdash \forall x[Q(x)]$$

You should only use rules from the following rule set, which was presented in the lecture slides, to construct your proof:

$$\frac{\text{Axiom}}{\alpha, \Gamma \vdash \alpha, \Delta}$$

$$\frac{\alpha, \beta, \Gamma \vdash \Delta}{(\alpha \wedge \beta), \Gamma \vdash \Delta} [\wedge\vdash] \quad \frac{\Gamma \vdash \alpha, \Delta \text{ and } \Gamma \vdash \beta, \Delta}{\Gamma \vdash (\alpha \wedge \beta), \Delta} [\wedge\vdash]$$

$$\frac{\alpha, \Gamma \vdash \Delta \text{ and } \beta, \Gamma \vdash \Delta}{(\alpha \vee \beta), \Gamma \vdash \Delta} [\vee\vdash] \quad \frac{\Gamma \vdash \alpha, \beta, \Delta}{\Gamma \vdash (\alpha \vee \beta), \Delta} [\vdash\vee]$$

$$\frac{\Gamma \vdash \alpha, \Delta}{\neg\alpha, \Gamma \vdash \Delta} [\neg\vdash] \quad \frac{\Gamma, \alpha \vdash \Delta}{\Gamma \vdash \neg\alpha, \Delta} [\vdash\neg]$$

$$\frac{\Gamma, \neg\alpha \vee \beta \vdash \alpha, \Delta}{\Gamma, \alpha \rightarrow \beta \vdash \Delta} [\rightarrow\vdash r.w.] \quad \frac{\Gamma \vdash \neg\alpha \vee \beta, \Delta}{\Gamma \vdash \alpha \rightarrow \beta, \Delta} [\vdash\rightarrow r.w.]$$

$$\frac{\forall x[\Phi(x)], \Phi(k), \Gamma \vdash \Delta}{\forall x[\Phi(x)], \Gamma \vdash \Delta} [\forall\vdash] \quad \frac{\Gamma \vdash \Phi(k), \Delta}{\Gamma \vdash \forall x[\Phi(x)], \Delta} [\vdash\forall]^\dagger$$

$^\dagger$  where  $\kappa$  cannot occur anywhere in the lower sequent.

**Answer:** The sequent is valid as shown by the following proof:

$$\begin{array}{c} \text{AXIOM} \\ \hline \text{Ax}[Px], Pa, \quad \text{Ax}[Px \rightarrow Qx], \quad |- \quad Pa, Qa \\ \hline \text{Ax}[Px], Pa, \quad \text{Ax}[Px \rightarrow Qx], \quad -Pa \quad |- \quad Qa \quad \quad \quad \text{AXIOM} \\ \hline \text{Ax}[Px], Pa, \text{Ax}[Px \rightarrow Qx], Qa \quad |- \quad Qa \\ \hline \text{Ax}[Px], Pa, \text{Ax}[Px \rightarrow Qx], -Pa \vee Qa \quad |- \quad Qa \\ \hline \text{Ax}[Px], Pa, \text{Ax}[Px \rightarrow Qx], Pa \rightarrow Qa \quad |- \quad Qa \quad \quad \quad [-\rightarrow \vdash r.w.] \\ \hline \text{Ax}[Px], Pa, \text{Ax}[Px \rightarrow Qx] \quad |- \quad Qa \quad \quad \quad [A \vdash -] \\ \hline \text{Ax}[Px], \text{Ax}[Px \rightarrow Qx] \quad |- \quad Qa \quad \quad \quad [A \vdash -] \\ \hline \text{Ax}[Px], \text{Ax}[Px \rightarrow Qx] \quad |- \quad Qa \quad \quad \quad [-\vdash A] \\ \hline \text{Ax}[Px], \text{Ax}[Px \rightarrow Qx] \quad |- \quad \text{Ax}[Qx] \end{array}$$

There is basically a mark for each correct rule application. Some credit may be given for almost correct rule applications.

**[Question 1 total: 20 marks]**

Answers

## Question 2

- (a) (i) Give the set of *clausal* formulae (i.e. formulae in *disjunctive normal form*) corresponding to the following propositional formulae: **[4 marks]**

$$\neg\neg A \vee S, (\neg S \wedge T), (A \vee B) \rightarrow Q, (Q \wedge T) \rightarrow (R \wedge S)$$

**Answer:**

$$\{ \{A, S\}, \{\neg S\}, \{T\}, \{\neg A, Q\}, \{\neg B, Q\}, \{\neg Q, \neg T, R\}, \{\neg Q, \neg T, S\} \}$$

- (ii) Give a proof that these formulae are inconsistent using *binary propositional resolution*. **[4 marks]**

<b>Answer:</b>	1. $\{A, S\}$	8. $\{A\}$	1&2
	2. $\{\neg S\}$	9. $\{Q\}$	4&8
	3. $\{T\}$	10. $\{\neg T, S\}$	7&9
	4. $\{\neg A, Q\}$	11. $\{\neg T\}$	2&10
	5. $\{\neg B, Q\}$	12. $\emptyset$	3&11
	6. $\{\neg Q, \neg T, R\}$		
	7. $\{\neg Q, \neg T, S\}$		

- (b) Translate the following sentence into *Propositional Tense Logic*: **[2 marks]**

If I win the lottery I will be rich forever after that.

**Answer:**  $G(Win \rightarrow GRich)$  (or  $G(PWin \rightarrow Rich)$ , though not quite right)

- (c) A Situation Calculus theory makes use of fluents of the forms:

$robot\_has(item)$        $on\_floor(item, room)$        $locked(door)$   
 $robot\_location(room)$        $connects(door, room_1, room_2)$

The theory includes constants referring to items, one of which is key.

The theory also describes the behaviour of a robot in terms of the following actions:

$pick\_up(object)$      $unlock(door)$      $move\_to(room)$

An initial situation,  $s_0$ , is described as follows:

Holds( $connects(door1, hall, lounge)$ ,  $s_0$ )    Holds( $connects(door2, hall, study)$ ,  $s_0$ )  
 $\neg$ Holds( $locked(door1)$ ,  $s_0$ )    Holds( $locked(door2)$ ,  $s_0$ )  
Holds( $on\_floor(key, lounge)$ ,  $s_0$ )    Holds( $robot\_location(hall)$ ,  $s_0$ )

- (i) Assuming that the initial situation is  $s_0$ , give a sequence of actions that will result in the goal  $robot\_location(study)$  being satisfied. **[2 marks]**
- (ii) For each of the actions  $pick\_up$  and  $move\_to$  specify a *precondition* axiom stating the conditions under which the action is possible. **[4 marks]**
- (iii) Give an *effect* axiom specifying the results of carrying out the action  $unlock$ . **[2 marks]**

- (iv) Write down a *frame* axiom stating that the `move_to` action does not affect the *locked* fluent.

**[2 marks]**

**Answer:**

- ```
(i) move(lounge), pickup(key), move(hall), unlock(door2), move(study)
(ii) - Poss(pickup(x),s) <- E r[ Holds(robot_location(r),s) &
                                     Holds(Onfloor(x,r),s)]
      - Poss(move(x),s) <- E r d[ Holds(connects(d,r,x),s) &
                                     Holds(robot_location(r),s) & - Holds(locked(d), s)]
```

(It is actually also ok to leave out the existential quantifier in these axioms, since the usual implicit universal quantification converts to existential when on right of  $<-$ .)

- (iii)  $-- \text{ Holds( locked(x), result(unlock(x),s)) } \leftarrow \text{ Poss(unlock(x),s)}$   
 (v)  $\text{ Holds(locked(x), result(move(x),s)) } \leftrightarrow \text{ Holds(locked(x),s)}$

**[Question 2 total: 20 marks]**

### Question 3

(a) For each of the following *Prolog* queries, give the value of the variable  $X$  after the query has been executed:

- (i) ?-  $X = 7/2$ . [1 mark]  
 (ii) ?-  $[1, [2, 3], 4] = [\_ | [X | \_ ]]$  [1 mark]  
 (iii) ?-  $A = [1, 2, 3, 4, 5], \text{setof}(I, (\text{member}(I, A), I > 2), X)$ . [1 mark]  
 (iv) ?-  $\text{append}([X], [2, 3], [1, 2, 3])$ . [1 mark]

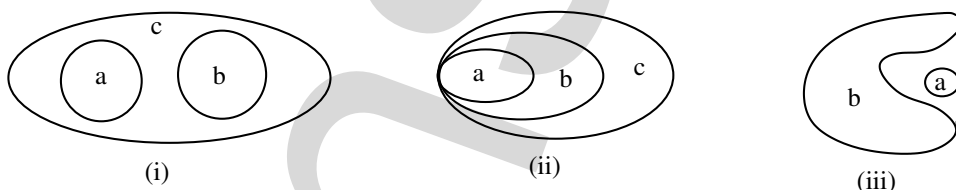
**Answer:**

- (i)  $7/2$   
 (ii)  $[2, 3]$   
 (iii)  $[3, 4, 5]$   
 (iv)  $1$

(b) Consider the following formulae involving topological relations of the *Region Connection Calculus* (RCC) and the *convex hull* function,  $\text{conv}$ . The constants ( $a$ ,  $b$  and  $c$ ) refer to particular spatial regions. In each case, draw a configuration of the regions referred to by these constants that satisfies the formula, labelling your diagram to indicate which region is which:

- (i)  $\text{DC}(a, b) \wedge \text{NTPP}(\text{sum}(a, b), c)$  [2 marks]  
 (ii)  $\text{TPP}(a, b) \wedge \text{TPP}(b, c) \wedge \text{TPP}(a, c)$  [2 marks]  
 (iii)  $\text{DC}(a, b) \wedge \text{TPP}(a, \text{conv}(b))$  [2 marks]

**Answer:** Possible diagrams are as follows:



(c) A *liger* is an animal whose parents are a male lion and a female tiger. Use *Description Logic* to give a definition of the concept **Liger** in terms of the concepts **Lion**, **Tiger**, **Male**, **Female** and the relation **hasParent**. [4 marks]

**Answer:**

$$\text{Liger} \equiv \exists \text{hasParent}.(\text{Male} \sqcap \text{Lion}) \sqcap \exists \text{hasParent}.(\text{Female} \sqcap \text{Tiger})$$

- (d) Write a *Default Logic* rule that formally represents the reasoning principle expressed in the following statement: **[2 marks]**

“British people typically drink tea, except for children and those who drink coffee.”

**Answer:**  $British(x) : Drinks(x, tea), \neg Child(x), \neg Drinks(x, coffee) / Drink(x, tea)$

- (e) This question concerns a *Fuzzy Logic* in which the following definitions of *linguistic modifiers* are specified:

$$\text{quite}(\phi) = \phi^{1/2} \quad \text{very}(\phi) = \phi^2$$

The logic is used to describe Leo the lion, who possesses certain characteristics to the following degrees:

$$\text{Large}(\text{leo}) = 0.5 \quad \text{Fierce}(\text{leo}) = 0.09 \quad \text{Clever}(\text{leo}) = 0.4$$

Translate the following sentences into fuzzy logic and also give the fuzzy truth value of each proposition (under the standard fuzzy interpretation of the Boolean connectives):

- (i) Leo is not very clever. **[2 marks]**  
(ii) Leo is very very large and quite fierce. **[2 marks]**

**Answer:**

A)  $\neg \text{very}(\text{Clever}(\text{leo}))$  (1 mark)

Truth value  $1 - (0.4)^2 = 1 - 0.16 = 0.84$  (1 mark)

B)  $\text{very}(\text{very}(\text{Large}(\text{leo}))) \wedge \text{quite}(\text{Fierce}(\text{leo}))$  (1 mark)

Truth value =  $\text{Min}(((0.5)^2)^2, (0.09)^{\frac{1}{2}}) = \text{Min}(0.0625, 0.3) = 0.0625$  (1 mark)

**[Question 3 total: 20 marks]**

**[Grand total: 60 marks]**